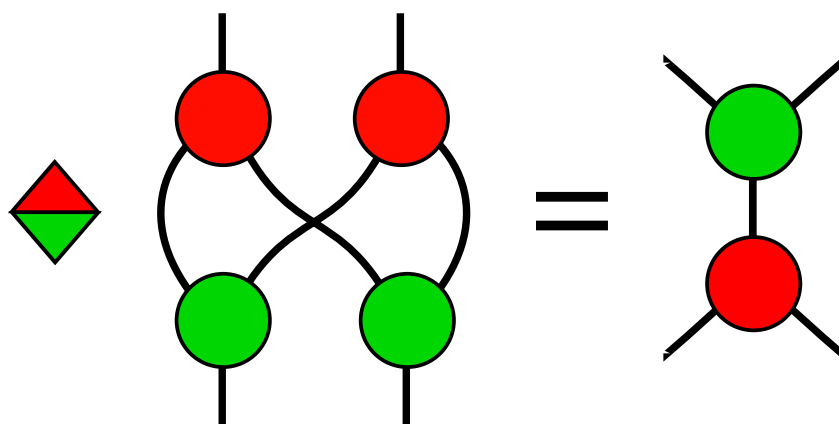


Mutually Unbiased Embeddings of Classical Logic



Ross Duncan

Oxford University Computing Lab

Joint work with Bob Coecke

Motivation

- Quantum observables may be incompatible: position/momentum, polarisation, spin ...
- In traditional quantum logic approaches these observables are simply *incomparable* in the lattice.
- However if one wants to *compute* with quantum mechanics we need know how these observables relate to each other.

We want to extract the *constructive content* from the incompatibility of quantum observables.

Classical Structures

Classical Structures capture the algebraic aspects of copying and deleting;

$$\delta = \text{green circle with one input and two outputs}$$

$$\epsilon = \text{green circle with one output and no inputs}$$

$$\delta^\dagger = \text{green circle with two inputs and one output}$$

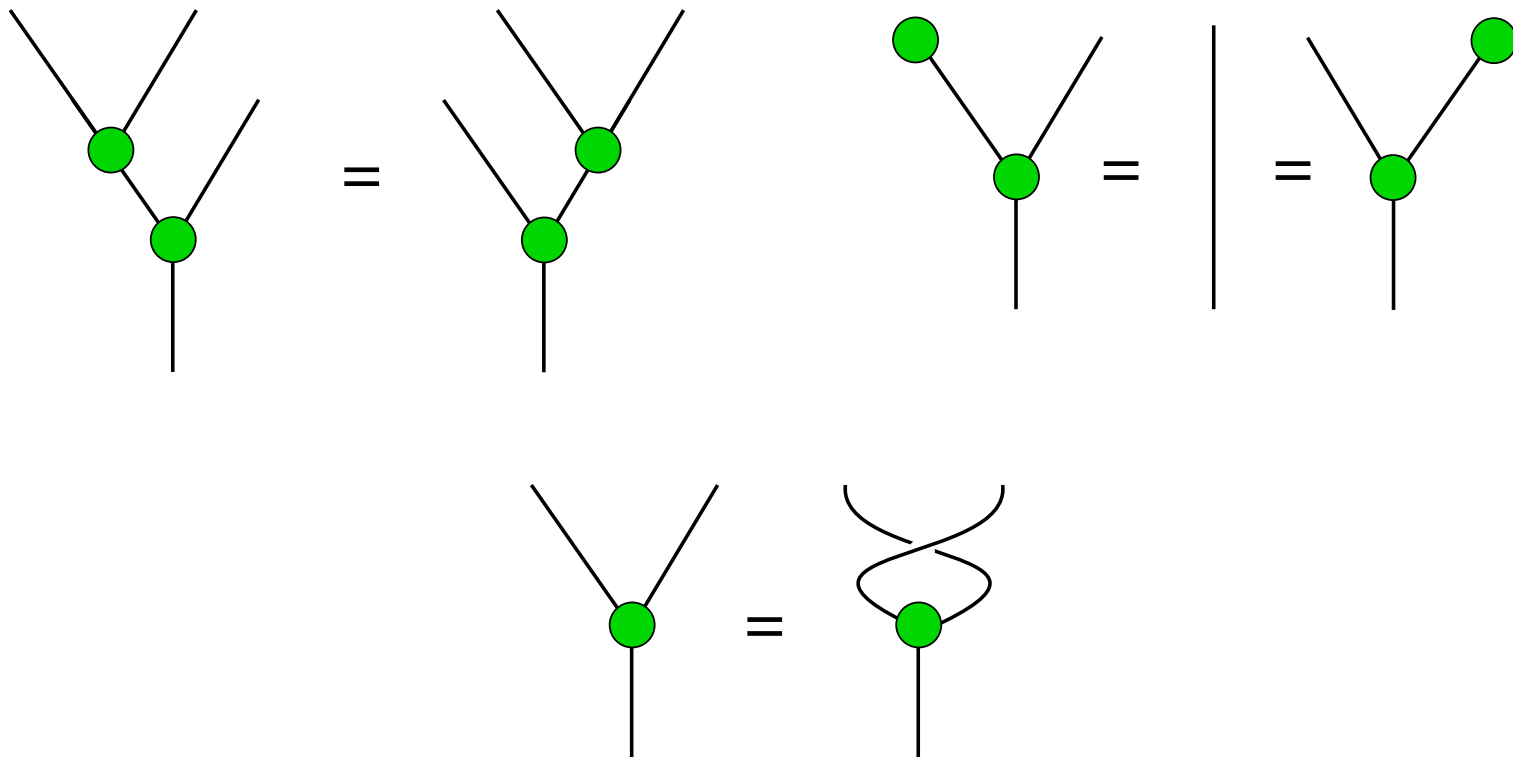
$$\epsilon^\dagger = \text{green circle with one input and no outputs}$$

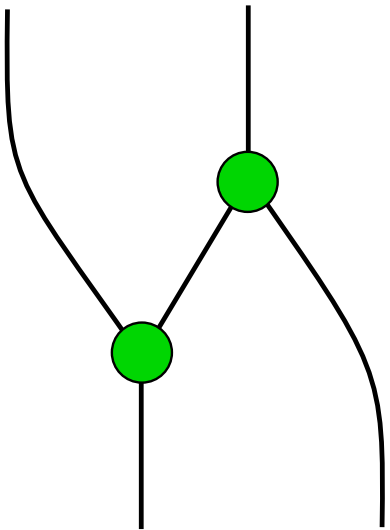
An example:

$$\delta : |i\rangle \mapsto |ii\rangle$$

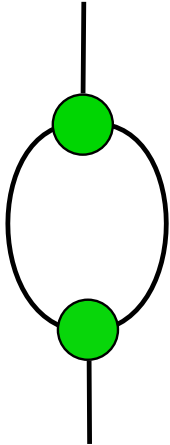
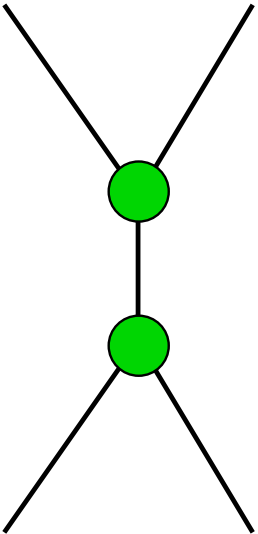
$$\epsilon : \sum_i |i\rangle \mapsto 1$$

Classical Structures





=



=



Cloning

Consider the map:

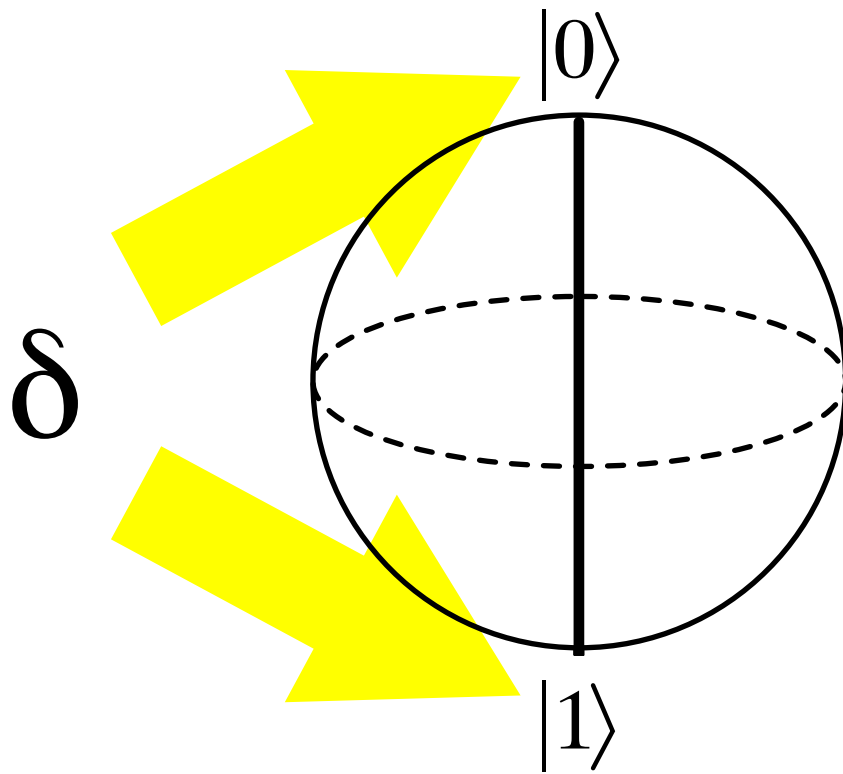
$$\delta_Z : Q \rightarrow Q \otimes Q :: |i\rangle \mapsto |ii\rangle$$

δ_Z is the *cloning* map for the basis $|0\rangle, |1\rangle$.

Obviously δ_Z is cannot clone all states:

$$\delta_Z |+\rangle = \delta_Z(|0\rangle + |1\rangle) = |00\rangle + |11\rangle$$

However, since quantum states are indistinguishable up to global phase the *vectors* $e^{i\alpha} |0\rangle$ and $e^{i\beta} |1\rangle$, are also cloned, when viewed as quantum states; hence can view δ as fixing an *observable* i.e. an axis of the Bloch sphere.



Deleting

Q: How to “erase” a quantum state $|\psi\rangle$ known to be in some given basis?

A: Use a measurement which gives *no information* about the existing state — i.e measurement in a basis $\{b_i\}$ such that

$$\begin{aligned} |\langle b_i | \psi \rangle| &= |\langle b_j | \psi \rangle| \\ \Rightarrow |\langle b_i | a_k \rangle| &= |\langle b_j | a_k \rangle| \\ \Rightarrow |\langle b_i | a_k \rangle| &= \frac{1}{\sqrt{d}} \text{ (in finite dim.)} \end{aligned}$$

Hence the idea of *Mutually Unbiased Bases* arise very naturally from the idea of *deleting* a classical value embedded in a quantum state space.

With respect to the Z -basis (i.e. $|0\rangle, |1\rangle$), the unbiased states are exactly:

$$\epsilon_Z^\alpha = (|0\rangle + e^{i\alpha} |1\rangle)/2$$

However if we compose ϵ_Z^α with δ_Z :

$$(\text{id} \otimes \epsilon_Z^\alpha) \circ \delta_Z = Z_{-\alpha} = \begin{pmatrix} 1 & 0 \\ 0 & e^{-i\alpha} \end{pmatrix}$$

Therefore the “copying” and “erasing” operations when combined induce a non-trivial action on states, unless $\alpha = 0$ (Will come back to this a bit later).

Thus, we have a classical structure:

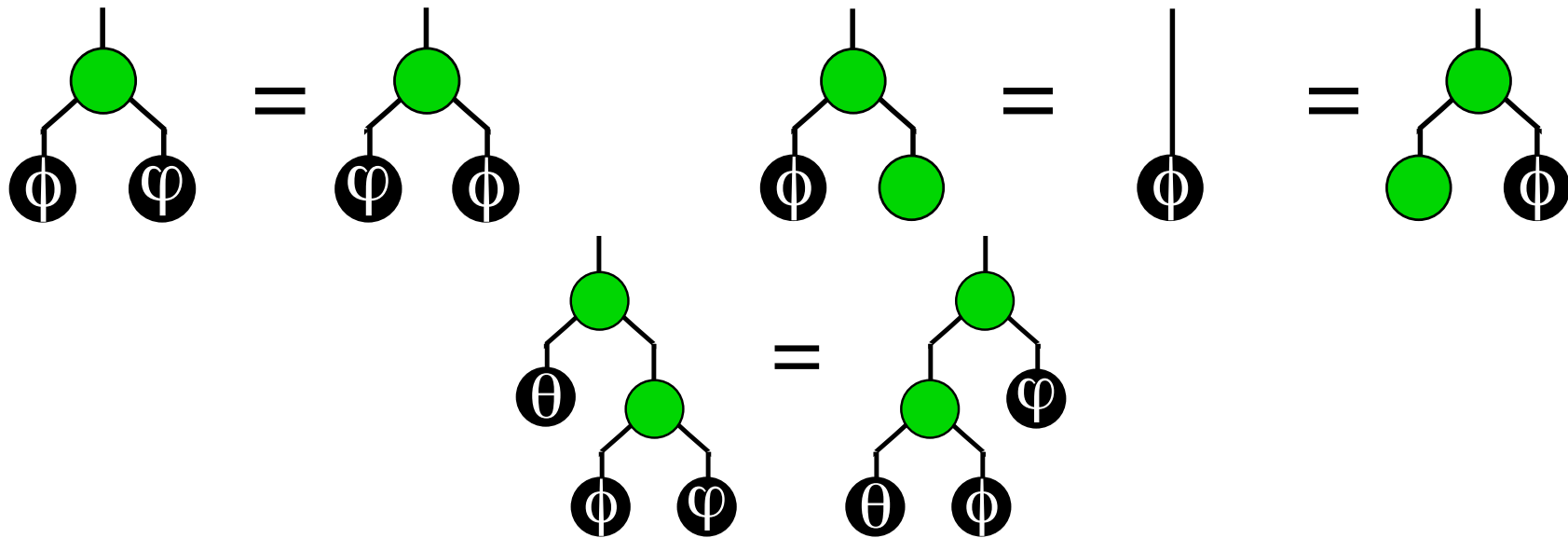
- δ_Z is the *cloning* map for the basis $|0\rangle, |1\rangle$.
- ϵ_Z is the *uniform deleting* of this basis.

The monoid of points

Let (A, δ, ϵ) be a classical structure; since $(\delta^\dagger, \epsilon^\dagger)$ is a commutative monoid we have a commutative multiplication on points

$\psi, \phi : I \rightarrow A$, defined by

$$\phi \odot \psi = \delta^\dagger \circ (\phi \otimes \psi)$$

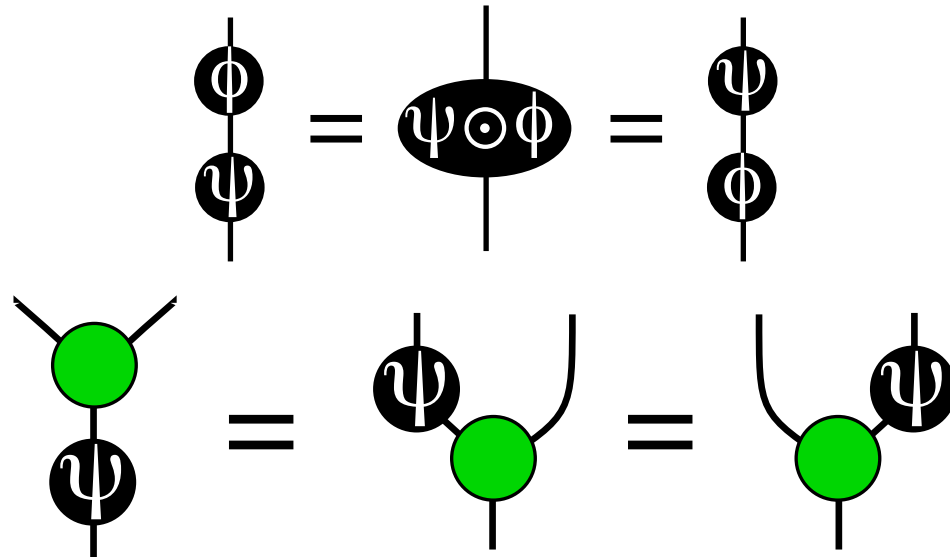


The monoid “currying”

Hence we have a lifting operation:

$$\Lambda : \mathcal{C}(I, A) \rightarrow \mathcal{C}(A, A) :: \psi \mapsto \delta^\dagger \circ (\psi \otimes \text{id}_A) \quad \Psi := \text{diagram}$$

This is homomorphism of monoids so:

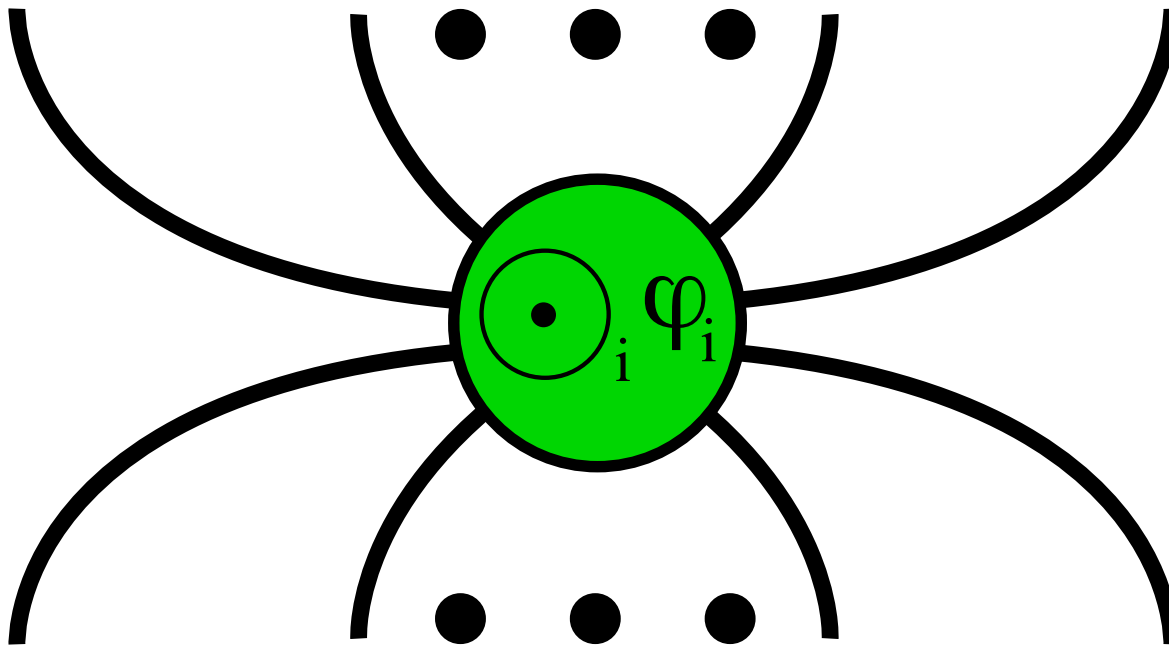


Also: $(\Lambda(\psi))^\dagger = \Lambda(\psi_*)$

Generalised Spider Theorem

Theorem. *Any map constructed from δ , ϵ , some points ψ_i and the adjoints of all of these, is determined only by its type and a product*

$\odot_i \psi_i$:



Unbiased Points

Proposition 1. *Let $\{|i\rangle\}_i$ be a basis for some finite dimensional Hilbert space, and let*

$$\delta : |i\rangle \mapsto |ii\rangle$$

be a cloning map for this basis; let $|\psi\rangle$ be any point. Then $\Lambda(\psi)$ is unitary if and only if $|\psi\rangle$ is unbiased with respect to $\{|i\rangle\}$.

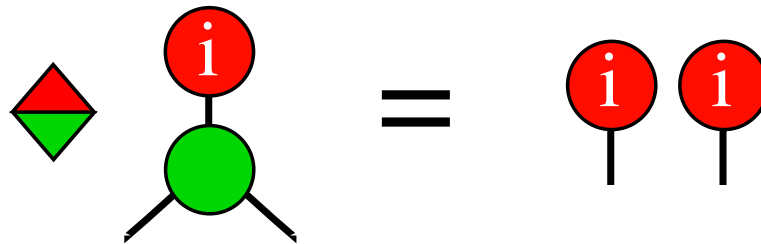
This allows us to *define* unbiasedness in the abstract setting.

Definition 2. A point $\psi : I \rightarrow A$ is unbiased for relative to (A, δ, ϵ) if $\Lambda(\psi)$ is unitary.

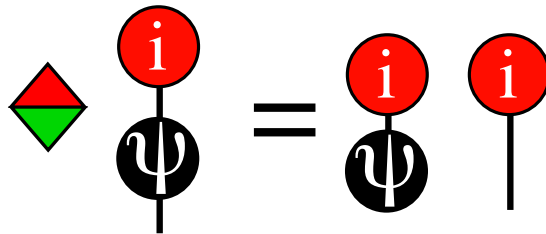
Note then that the set of unbiased points form an abelian group with respect to the \odot operation.

Classical Points

Definition 3. A point $i : I \rightarrow A$ is called *classical* relative to (A, δ, ϵ) if $\delta \circ i = i \otimes i$.



The classical points are eigenvectors in a suitable sense:



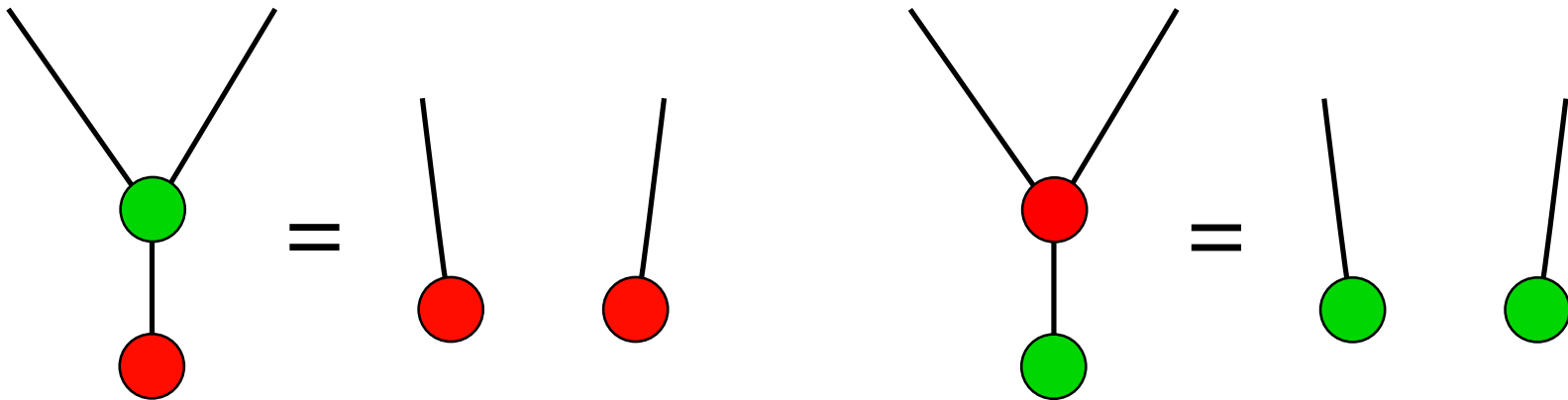
Complementary Classical Structures

Definition 4. *Two classical structures $(A, \delta_X, \epsilon_X)$ and $(A, \delta_Z, \epsilon_Z)$ in a \dagger -SMC are called complementary if they obey the following rules:*

- *whenever $z_i : I \rightarrow A$ is classical for (δ_X, ϵ_X) it is unbiased for (δ_Z, ϵ_Z) ;*
- *whenever $x_j : I \rightarrow A$ is classical for (δ_Z, ϵ_Z) it is unbiased for (δ_X, ϵ_X) ;*
- *ϵ_X^\dagger is classical for (δ_Z, ϵ_Z) and ϵ_Z^\dagger is classical for (δ_X, ϵ_X) .*

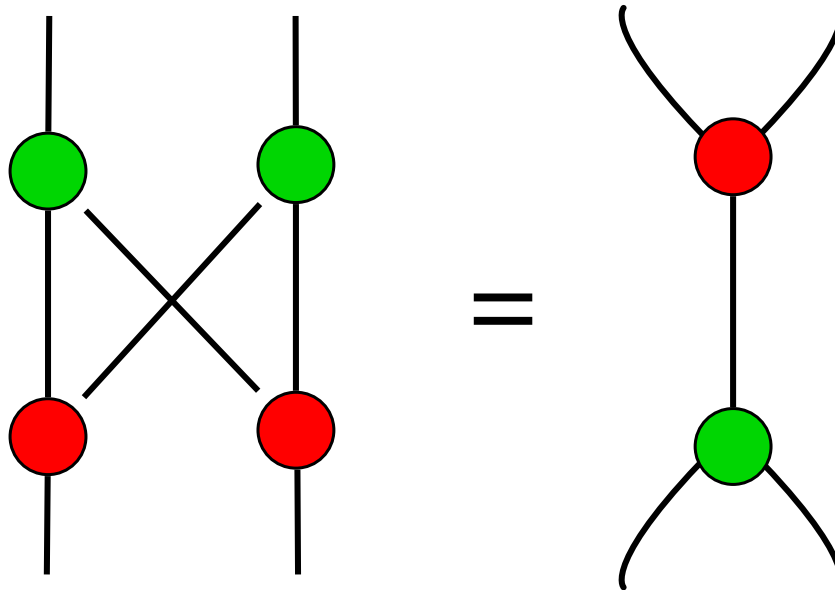
Bialgebraic Laws for Complementary Classical Structures

Cloning Laws:



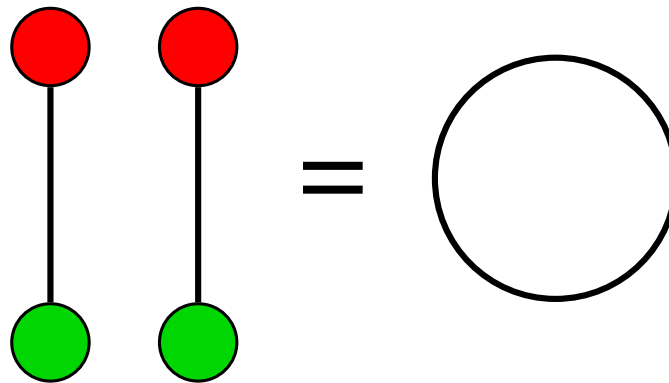
Bialgebraic Laws for Complementary Classical Structures

Bialgebra Law:



Bialgebraic Laws for Complementary Classical Structures

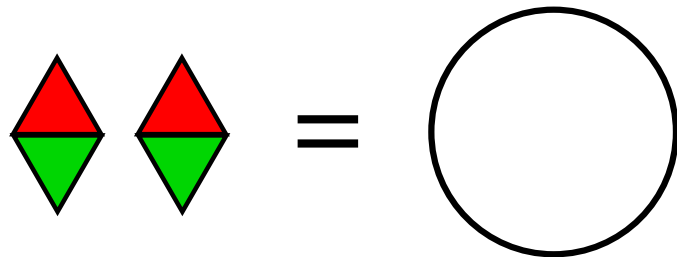
Dimension Law:



The pair of non-commuting observables fails to be a true bialgebra: every equation has a (hidden) scalar factor. Call this structure a *scaled bialgebra*.

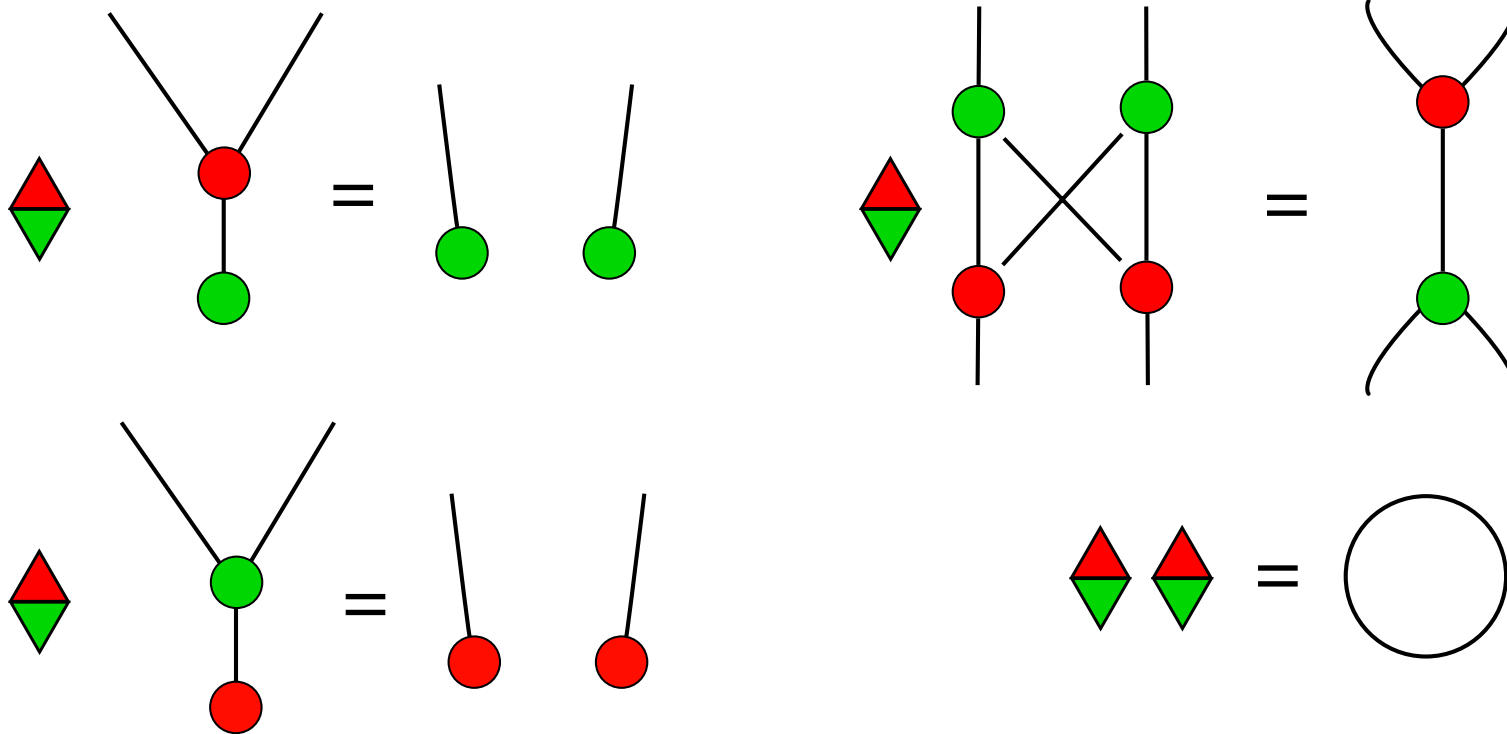
Bialgebraic Laws for Complementary Classical Structures

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Scaled Bialgebra Laws

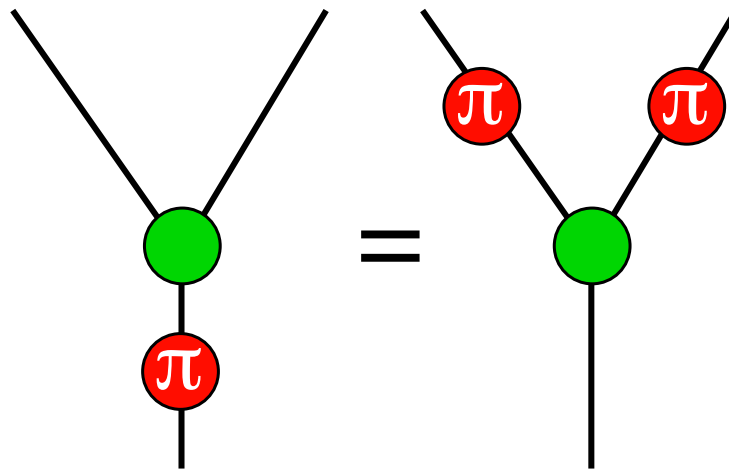


“Negation”

$$X_\pi = X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \because \begin{cases} |0\rangle \mapsto |1\rangle \\ |1\rangle \mapsto |0\rangle \end{cases}$$

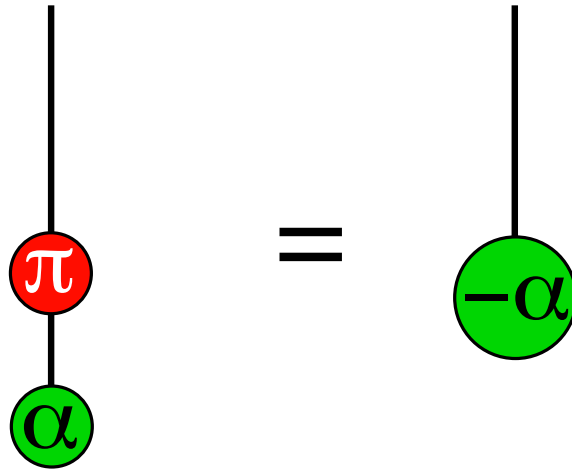
$$\begin{array}{ccc} Q & \xrightarrow{\delta} & Q \otimes Q \\ \downarrow X & & \downarrow X \otimes X \\ Q & \xrightarrow{\delta} & Q \otimes Q \end{array}$$

“Negation”



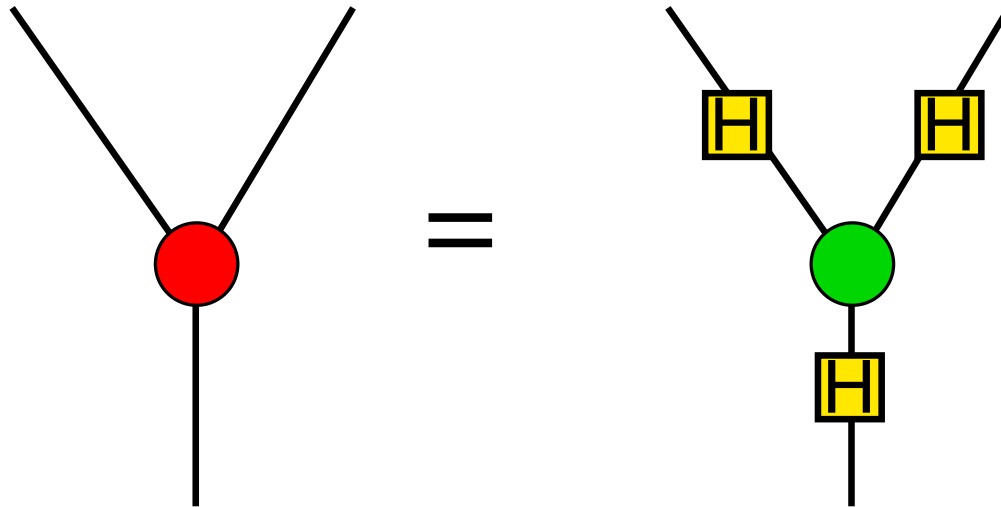
“Negation”

$$X :: |0\rangle + e^{i\alpha} |1\rangle \mapsto e^{i\alpha} |1\rangle + |0\rangle = |0\rangle + e^{-i\alpha} |1\rangle$$



Hadamard as a Mediating Map

We can define the red classical structure in terms of H and the green structure:

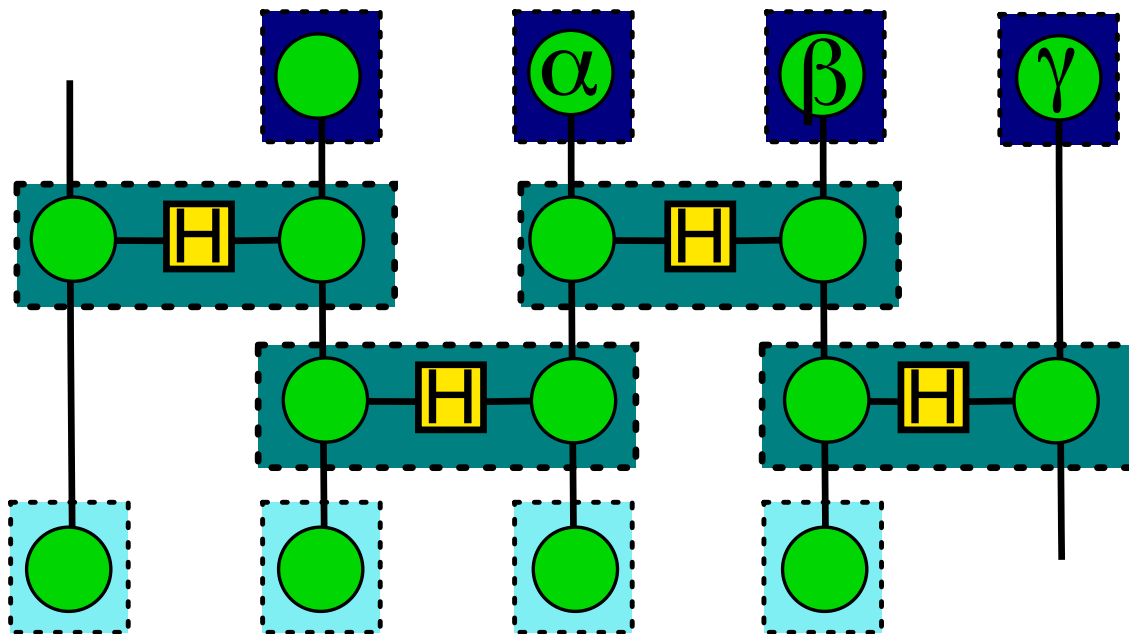


We can immediately derive a law for changing the colour of dots by introducing H boxes – in fact this gives a general “colour duality”.

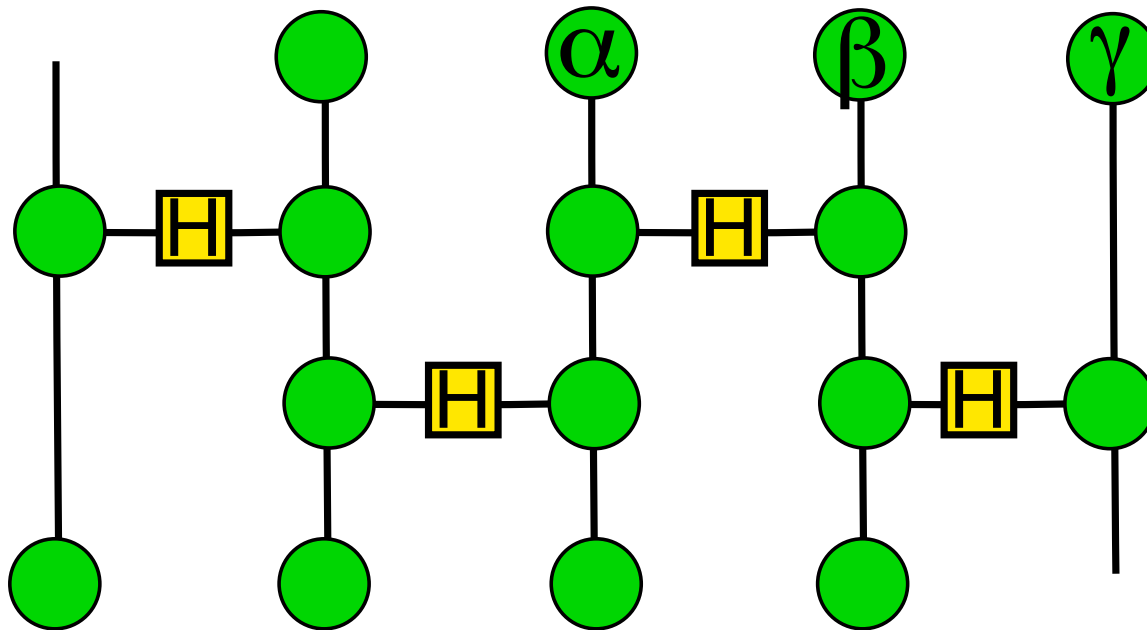
Example: A MBQC 1-qubit unitary

Proposition 5. *If U is a unitary on \mathbb{C}^2 there exist α, β, γ such that $U = Z_\alpha X_\beta Z_\gamma$.*

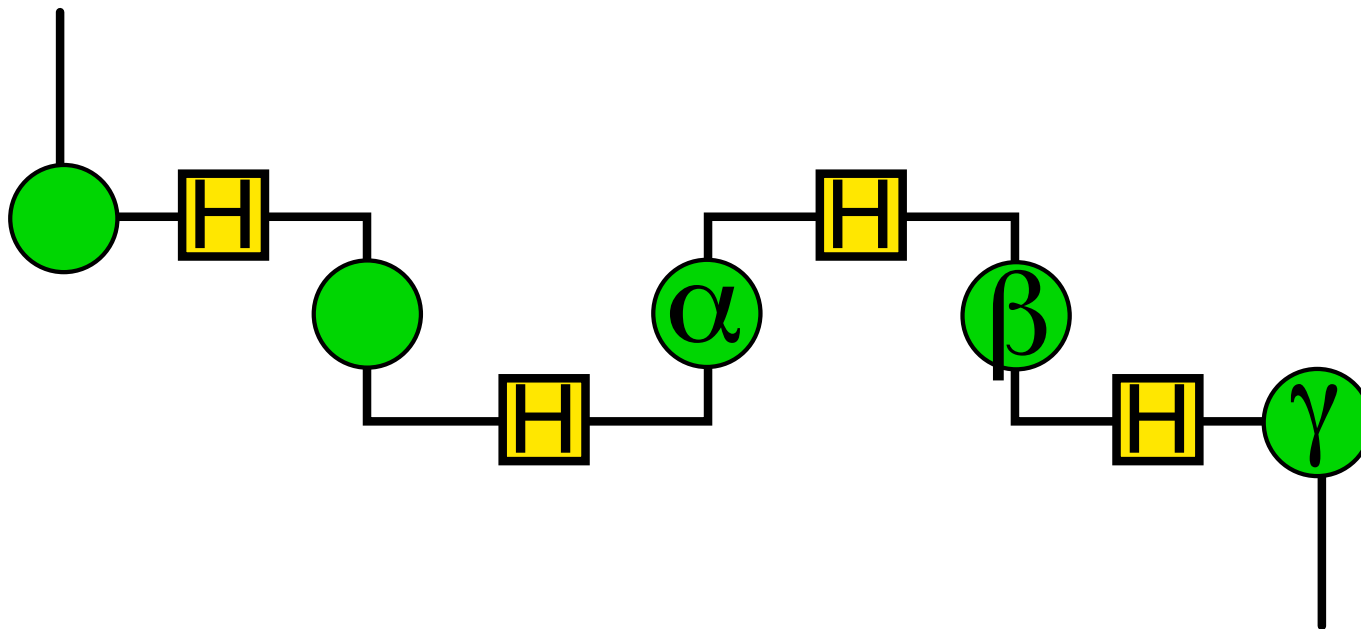
Here is (part of) a measurement based program to compute this:



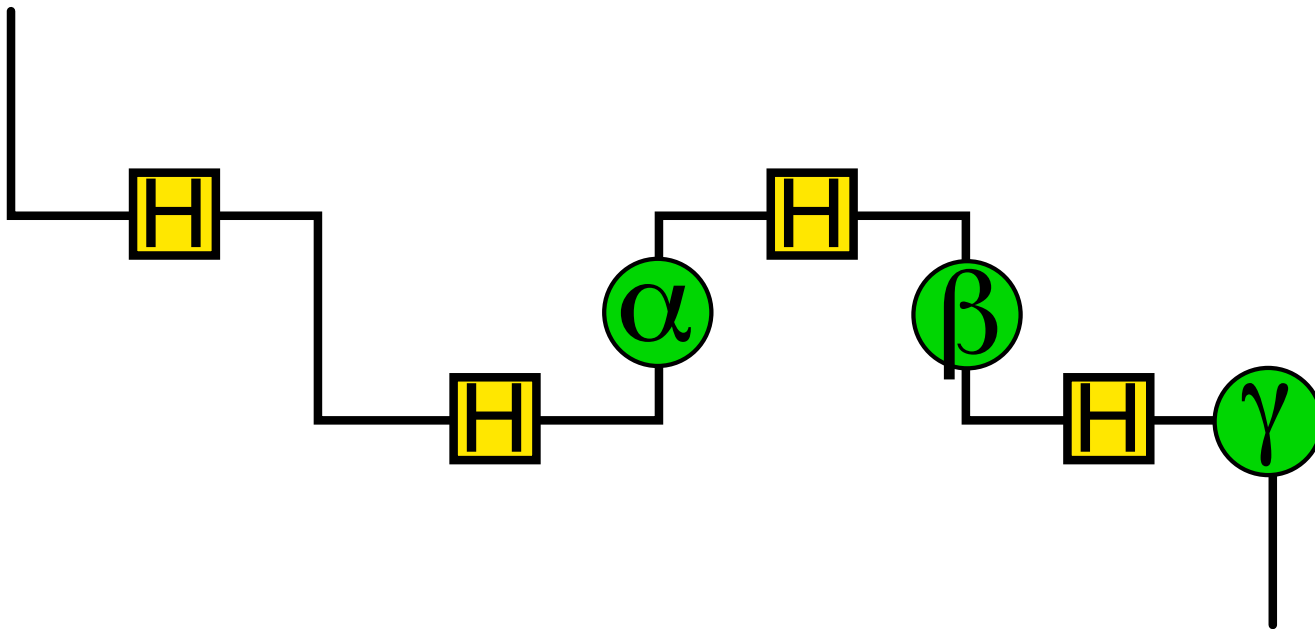
Example: A MBQC 1-qubit unitary



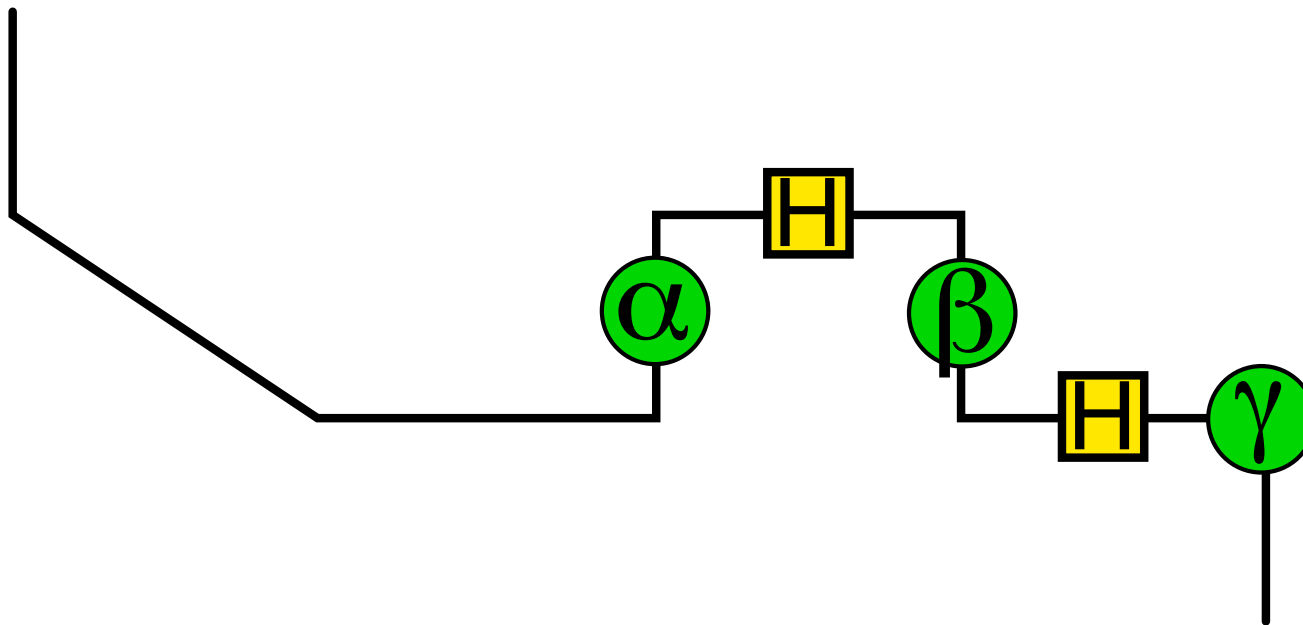
Example: A MBQC 1-qubit unitary



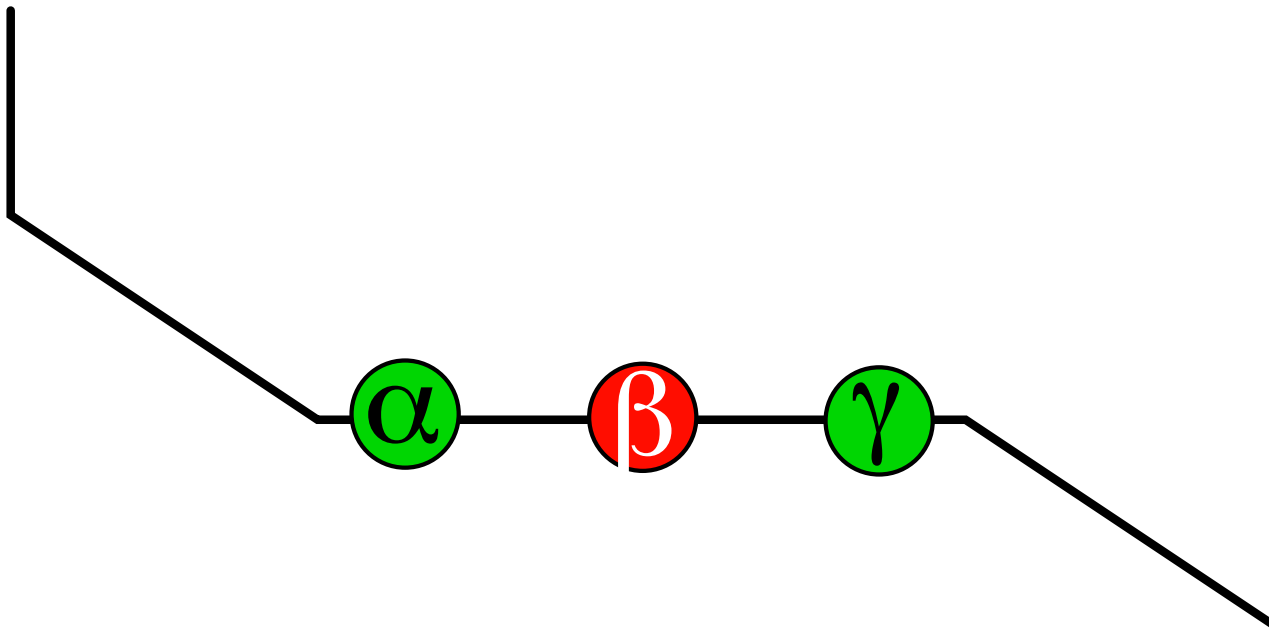
Example: A MBQC 1-qubit unitary



Example: A MBQC 1-qubit unitary



Example: A MBQC 1-qubit unitary



$$= Z_{\alpha} X_{\beta} Z_{\gamma}$$

Example: Quantum Fourier Transform

Among the most important quantum algorithms, the quantum Fourier transform is a key stage of factoring.

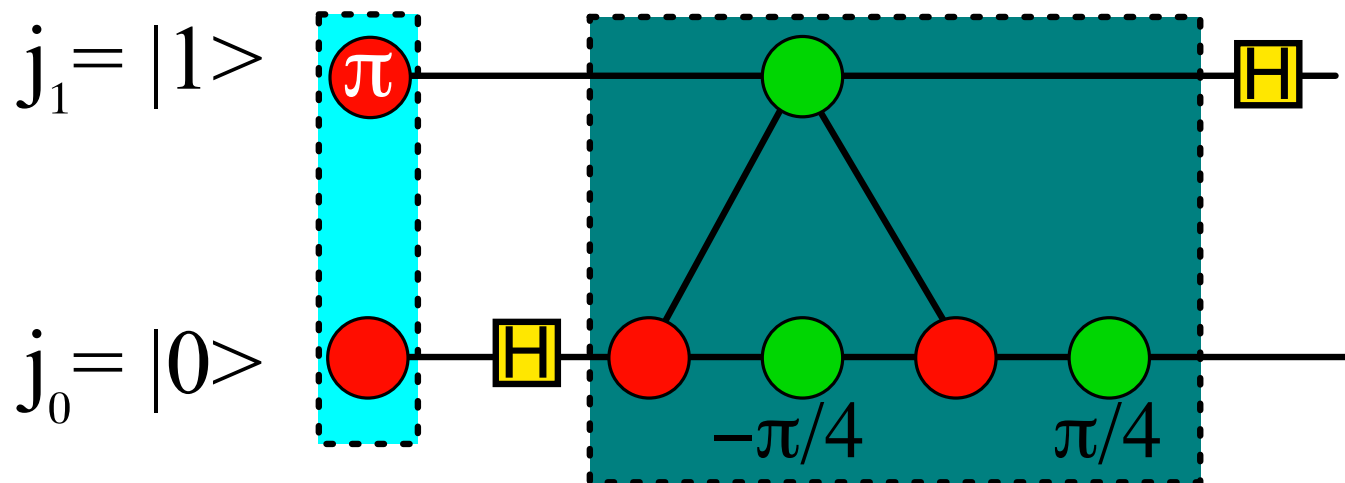
$$|j_0 j_1 \cdots j_n\rangle \mapsto (|0\rangle + e^{2\pi i \alpha_0} |1\rangle)(|0\rangle + e^{2\pi i \alpha_1} |1\rangle) \cdots (|0\rangle + e^{2\pi i \alpha_n} |1\rangle)$$

where $\alpha_k = 0.j_k \cdots j_n = \sum_{l=k}^n j_l / 2^k$

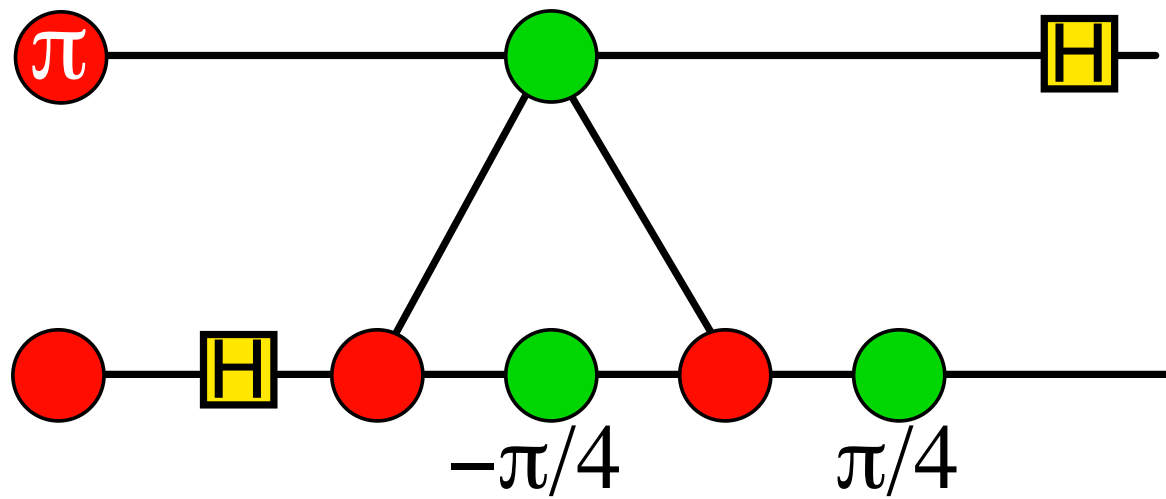
For 2 qubits:

$$\begin{aligned} |00\rangle &\mapsto (|0\rangle + |1\rangle)(|0\rangle + |1\rangle) & |10\rangle &\mapsto (|0\rangle + e^{i\pi} |1\rangle)(|0\rangle + |1\rangle) \\ |01\rangle &\mapsto (|0\rangle + e^{i\pi/2} |1\rangle)(|0\rangle + e^{i\pi} |1\rangle) & |11\rangle &\mapsto (|0\rangle + e^{i3\pi/2} |1\rangle)(|0\rangle + e^{i\pi} |1\rangle) \end{aligned}$$

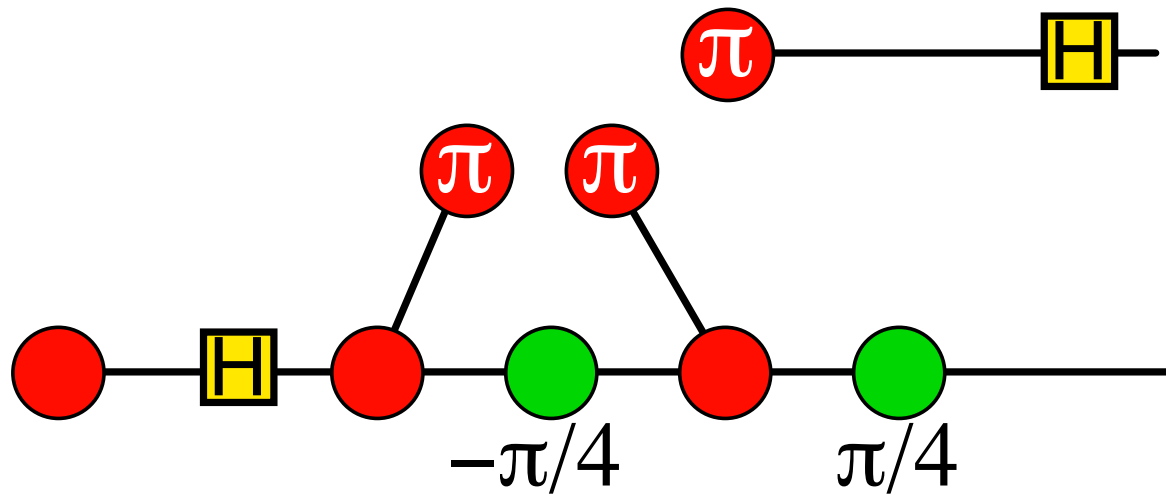
Example: Quantum Fourier Transform



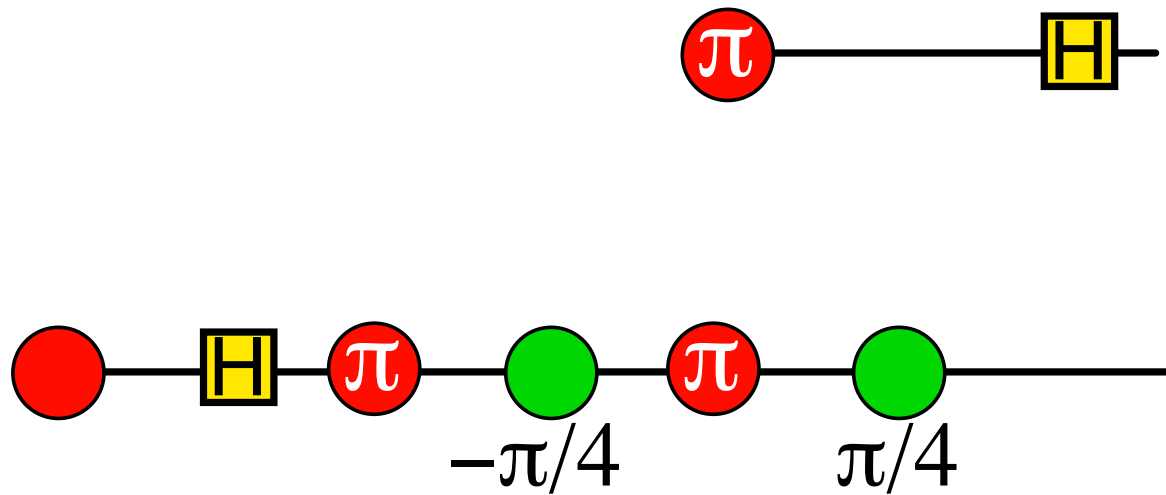
Example: Quantum Fourier Transform



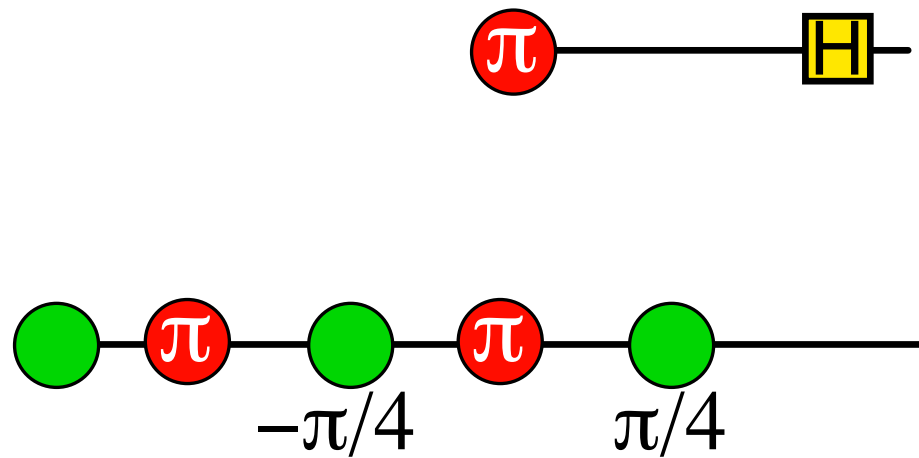
Example: Quantum Fourier Transform



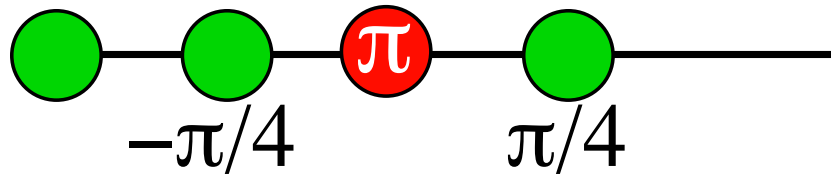
Example: Quantum Fourier Transform



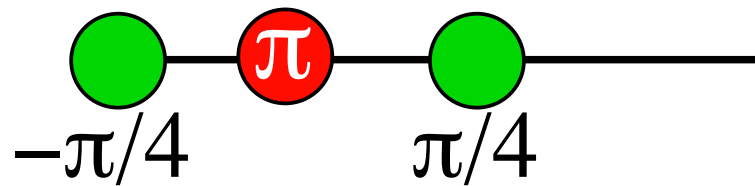
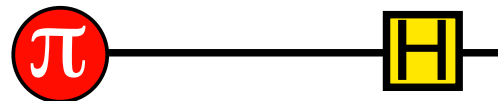
Example: Quantum Fourier Transform



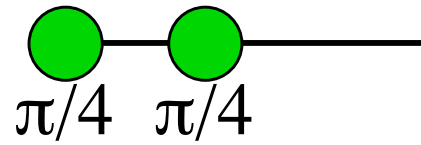
Example: Quantum Fourier Transform



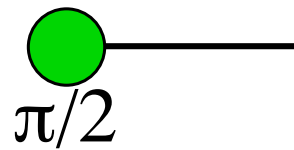
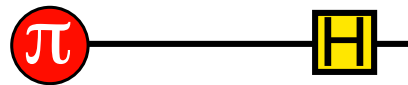
Example: Quantum Fourier Transform



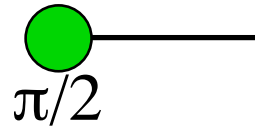
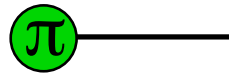
Example: Quantum Fourier Transform



Example: Quantum Fourier Transform



Example: Quantum Fourier Transform



which is the correct result! YAY!

To hear more:

Clifford Lectures at Tulane University this week, starting wednesday.

Information flow in physics, geometry, logic and computation

On saturday:

- 9.00: Samson Abramsky, *part 4*
- 10.30: Ross Duncan, *Logic of Complementary Quantum Observables*
- 14.00: Bob Coecke, *Kindergarten Quantum Mechanics*
- 15.30: Samson Abramksy, *part 5*

<http://www.math.tulane.edu/mwm/clifford>